

# CS551

## Layering and Addressing

### Bill Cheng

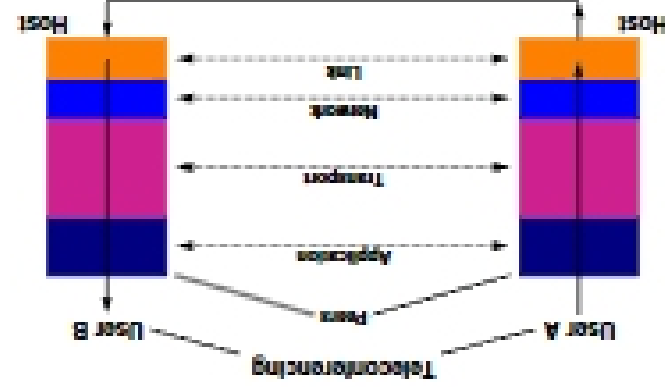
<http://merlot.usc.edu/cs551-f12>



- Set of rules governing communication between network elements (applications, hosts, routers)
- Protocols define:
  - = Format and order of messages
  - = Actions taken on receipt of a message
- Protocols are hard to design
  - = We need design guidelines!



- Layering: technique to simplify complex systems



## Layering Characteristics

- Each layer relies on services from layer below and exports services to layer above
- Interface defines interaction
- Hides implementation - layers can change without disturbing other layers (black box)



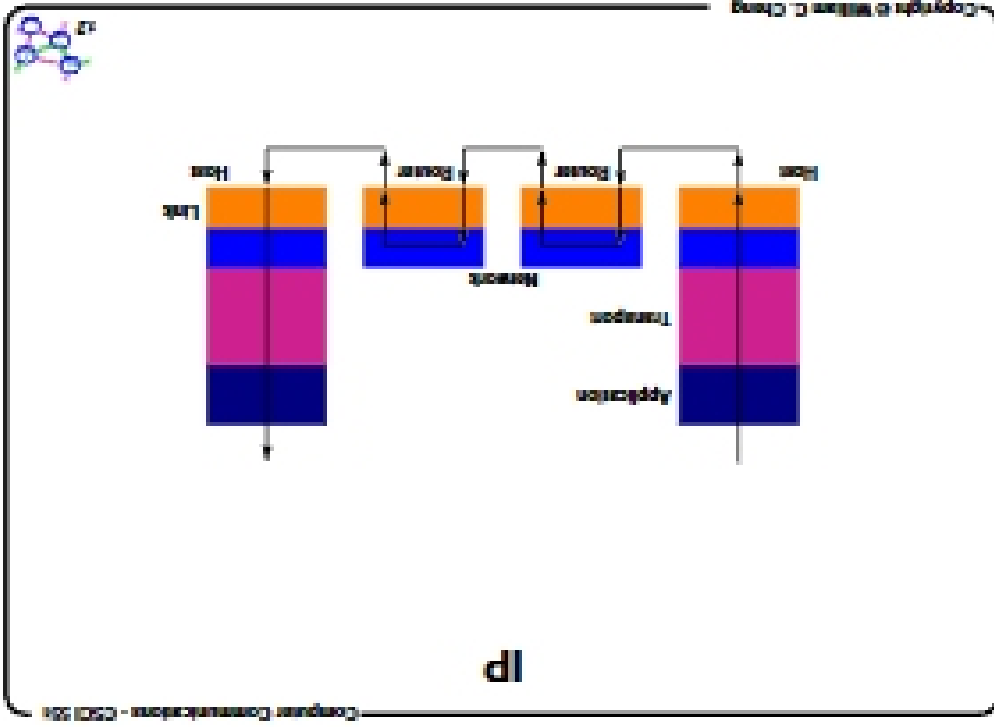
## Layer Encapsulation



## OSI Model: 7 Protocol Layers

- Physical: how to transmit bits
- Data link: how to transmit frames
- Network: how to route packets hop/hop
- Transport: how to send packets end/end
- Session: how to tie flows together
- Presentation: byte ordering, security
- Application: everything else!





- ### Is Layering Harmful?
- ◀ Sometimes,
    - ◀ Layer N may duplicate lower level functionality (e.g., error recovery).
    - ◀ Layers may need same info (timestamp, MTU).
    - ◀ Strict adherence to layering may hurt performance.
    - ◀ Naive layer implementation frequently hurts performance.

- ### Example: Transport Layer
- ◀ First end-to-end layer
  - ◀ End-to-end state
  - ◀ May provide reliability, flow control, and congestion control

# IP & TCP

## Course Focus

- ### Example: Network Layer
- ◀ Point-to-point communication
  - ◀ Network and host addressing
  - ◀ Routing

- ### Layering General Issues
- ◀ Reliability
  - ◀ Flow control
  - ◀ Fragmentation
  - ◀ Multiplexing
  - ◀ Connection setup (handshaking)
  - ◀ Addressing/naming (locating peers)

Copyright © William C. Stevens

### IP Header

Example Internet Datagram Header

```

0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
|-----|-----|-----|-----|-----|-----|-----|-----|
| Version | IHL | Type of Service | Total Length | Identification |
| Flags | Fragment Offset | Header Checksum | Time to Live | Protocol |
| Header Checksum | Source IP Address | Destination IP Address | | | |
|---|---|---|---|---|---|
| Options | Padding |
|-----|-----|-----|-----|-----|-----|

```

Copyright © William C. Stevens

### Fragmentation

- Forwarding costs per packet
- Nice if we can send large chunks of data
- Different link-layers have different MTUs
- Fragmentation
- Intra-network
- Inter-network

Copyright © William C. Stevens

### Path MTU Discovery

- Hosts dynamically discover MTU of path
- Send message with don't fragment bit
- Get ICMP message indicating size
- What happens if path changes?
- Increasing/decreasing path MTU
- Usually implemented by the *transport* layer
- Expected that future routing protocols will provide MTU information

Copyright © William C. Stevens

### IP Functions

- Type of service
- Not used until recently
- Identification, flags and fragment offset
- Fragmentation
- Time to live
- Bounded delivery
- Protocol
- (De)multiplexing higher layer protocols (analogous to *port numbers* in TCP)
- Length
- IP packet length limited to 64K
- Header checksum
- Ensures some degree of header integrity

Copyright © William C. Stevens

### Fragmentation Is Harmful

- Uses resources poorly
- Example of packet just bigger than MTU
- Poor end-to-end performance
- Loss of a fragment
- Reassembly is hard
- Buffering constrains

Copyright © William C. Stevens

### Path MTU

Algorithm:

- Initialize MTU to MTU to next hop
- Send datagrams with DF bit set
- If "datagram too big", decrease MTU
- Periodically (>5mins, or >1min after previous increase), increase MTU
- Some routers will return proper MTU
- MTU values cached in routing table