

key

5,  $\bar{2}$ , 4, 6, 1, 3

$\bar{2}$ , 5, 4, 6, 1, 3

2, 5,  $\bar{4}$ , 6, 1, 3

2,  $\bar{4}$ , 5, 6, 1, 3

2, 4, 5,  $\bar{6}$ , 1, 3

2, 4, 5, 6,  $\bar{1}$ , 3

2, 4, 5,  $\bar{1}$ , 6, 3

2, 4,  $\bar{1}$ , 5, 6, 3

2,  $\bar{1}$ , 4, 5, 6, 3

$\bar{1}$ , 2, 4, 5, 6, 3

1, 2, 4, 5, 6,  $\bar{3}$

1, 2, 4, 5,  $\bar{3}$ , 6

1, 2, 4,  $\bar{3}$ , 5, 6

1, 2,  $\bar{3}$ , 4, 5, 6



In the key outside of array.

# How to calculate running time?

- Each “line” of **pseudocode** requires a constant time. (In **RAM** model)

for  $j \leftarrow 2$  to  $length[A]$

do begin

$key \leftarrow A[j];$

$i \leftarrow j - 1;$

while  $i > 0$  and  $A[i] > key$

do begin

$A[i + 1] \leftarrow A[i];$

$i \leftarrow i - 1;$

end - while;

$A[i + 1] \leftarrow key;$

end - for.

This loop runs  $n-1$  times  
and each time runs at most  
 $4+3(j-1)$  lines.

This loop runs at most  
 $j-1$  times and each time  
runs at most 3 lines.

$$T(n) \leq \sum_{j=2}^n (4 + 3(j-1))$$
$$= \cancel{n} - 1 + 3 \frac{(n-1)(n+2)}{2}$$
$$= 4(n-1) + \frac{3(n-1)(n+2)}{2}$$