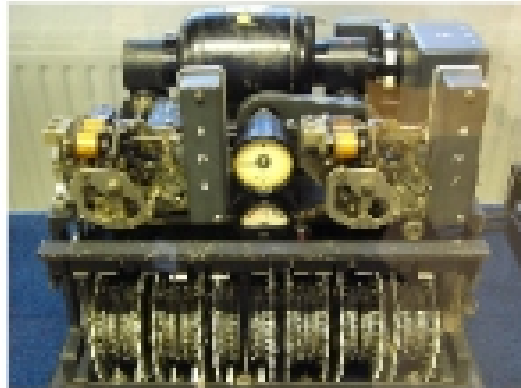


Class 25: Security through Complexity?



Lorenz cipher used in WWI I

FS6 is due today.

Karsten Nohl

CS439: Theory of Computation
University of Virginia, Computer Science

Motivation

- Many applications require certain tasks to be *easy* for some and *hard* for others
- Example: Decryption of encrypted message is *easy* only when given a secret key

Cryptography is concerned with constructing algorithms that withstand abuse. -Bellare&Shoup

Complexity is a powerful tool to "lock out" adversaries.
Basic Idea: Require *hard* problem to be solved, give hint as key.

Lecture 25: Complexity - Lecture 1

2

Computer Science

NP can be useful

- So far, you learnt how to detect "unsolvable" problems (in NP) and solve them anyway by approximation (in P)
- For cryptography we want the opposite: problems that are almost always *hard*, i.e., cannot be approximated in P

Lecture 25: Complexity - Lecture 1

3

Computer Science

(Breaking a strong cipher should require)
"as much work as solving a system of
simultaneous equations in a large number
of unknowns of a complex type"
- Shannon, '49

Sounds NP-Complete,
doesn't it?



Lecture 25: Complexity - Lecture 1

4

Computer Science

Goal: Encryption

- For almost all security schemes we need:
 - Encryption / one-way function

easy to compute: $enc(x, k) \rightarrow y$

hard to find any part of: $(x, k) = enc^{-1}(y)$

- Often also required:

- Decryption

$dec(y, k) \rightarrow x$

Lecture 25: Complexity - Lecture 1

5

Computer Science

Encryption build on Hardness

- Knapsack problem is NP-Complete
 - Problem of filling bag with best selection of items
 - Recall: Reducible from Subset-Sum
- Enable Encryption: Keep message secret by hiding it in a Knapsack instance

bits of encryption key = knapsack instance

$$S = \sum_{i=1}^n x_i a_i$$

message bits

Decryption possible by knowing easy knapsack instance (secret key) that provides shortcut.

Lecture 25: Complexity - Lecture 1

6

Computer Science

Flawed Security Argument

- Subset Sum is NP-Complete
- Breaking knapsack cipher involves solving a subset sum problem
- Therefore, knapsack cipher is secure

Flaw: NP-Complete means there is no **fast general solution**. Some instances may be solved quickly.

[Note: Adi Shamir broke knapsack cipher (1982)]

Cipher Design

- NP-Completeness is not sufficient for cryptographic hardness
Worst-case complexity
- Need solution to usually be hard
Average complexity
- Captured in new complexity class:
All *tractable* problems are in BPP
(which only makes sense if $P \neq NP$)

probabilistic:
can flip coins

Cipher Design (cont.)

- A “strong” cipher cannot be broken faster than exhaustive key search (brute force)

$\Theta(2^n)$ time

- Only possible shortcut:
Trade space for time; e.g.:

$\Theta(2^{n-5})$ time + space

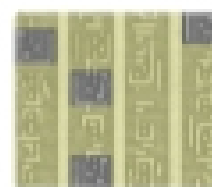
Results of Insufficient Hardness

- All broken cipher have a gap between *worst-case* and *average* hardness
- Estimating average hardness is often impossible (= finding best algorithm for instances of NP-complete problem)
- Next: Analyze cipher, identify complexity, and break it by finding tractable average solution.

Proprietary Cryptography (or: why “security-by-obscurity” never works)

First: Disclosure

- Secret algorithm can often be found:
 - Disassembling software
 - Hardware reverse-engineering



This talk: Breaking a cipher once we found it.

Then: Exploitation

- Most secret ciphers are broken after disclosure
- Flaws are very similar in all DIY ciphers (and cryptanalyst spot them in a glimpse)

"No more weak ciphers. No more paranoia."
Sean O'Neil

The crux of most flaws

- Most weaknesses caused by insufficient *non-linearity*.
- At the heart of the problem:
LFSRs (linear feedback shift register)



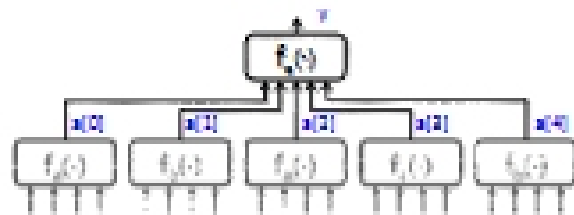
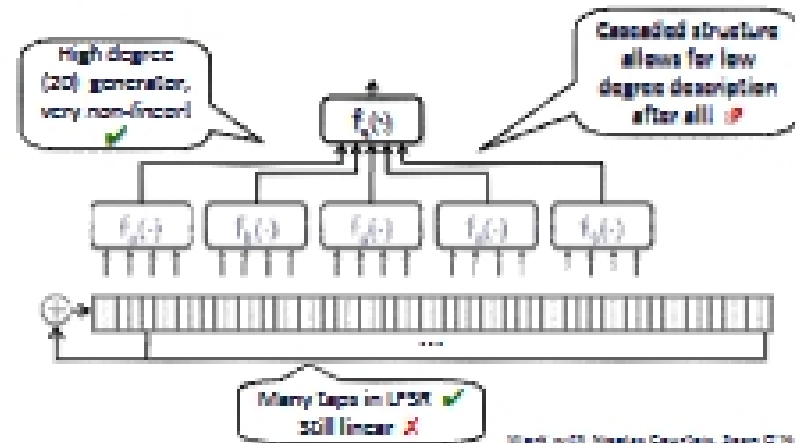
```
tmp = x[12]^x[15]^x[16]^x[17];
for i=17:-1:1  x[i]=x[i-1];
x[0] = tmp;
```

Non-Linearity

- System of equations that describes n -bit cipher can have up to $O(2^n)$ terms.
- Only $O(n)$ of these terms are linear.

Linear * P
Non-linear * NP

Mifare Crypto-1



Compute equations for first output bit:

```
a[0] = f(a[x[7], x[9], x[11], x[13]]);
a[1] = ...
...
y = f(a[0], a[1], a[2], a[3], a[4])
```

Before computing next bit, shift LFSR:

```
tmp = x[0]^...^x[17];
for i=1:17  x[i]=x[i+1];
x[18] = tmp;
```

Describes cipher as system of equations with $48+25$ unknowns, terms with degree $\leq 4!$

Almost there ...

1. Describe weak parts of cipher as system of equations
2. Brute-Force through complex parts:
Guess-and-Determine attack.
3. Solve system of equations:
MiniSAT is our friend



Solving for 48-bit Crypto-1 key takes 12 seconds; compared to month for brute-force.