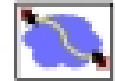




## 15-441 Computer Networking

### Lecture 18 – TCP Performance

## Outline



- TCP congestion avoidance
- TCP slow start
- TCP modeling

11-01-07

Lecture 18: TCP Congestion Control

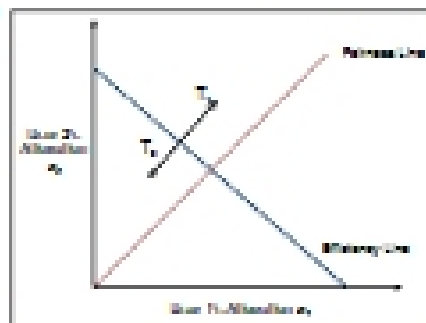
2

## Additive Increase/Decrease



- Both  $X_1$  and  $X_2$  increase/ decrease by the same amount over time

- Additive increase improves fairness and additive decrease reduces fairness



11-01-07

Lecture 18: TCP Congestion Control

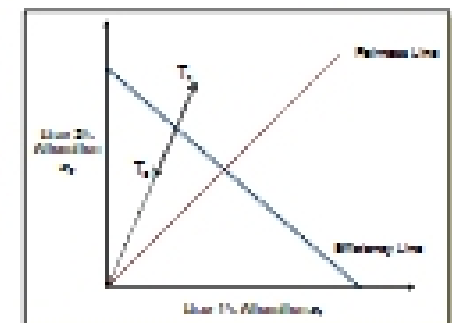
3

## Multiplicative Increase/Decrease



- Both  $X_1$  and  $X_2$  increase by the same factor over time

- Extension from origin – constant fairness

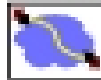


11-01-07

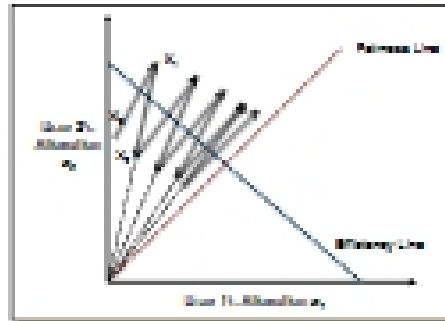
Lecture 18: TCP Congestion Control

4

## What is the Right Choice?



- Constraints limit us to AIMD
  - Improves or keeps fairness constant at each step
  - AIMD moves towards optimal point



11-0-07

Lecture 10: TCP Congestion Control

6

## TCP Congestion Control



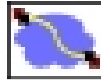
- Changes to TCP motivated by ARPANET congestion collapse
- Basic principles
  - AIMD
  - Packet conservation
  - Reaching steady state quickly
  - ACK clocking

11-0-07

Lecture 10: TCP Congestion Control

7

## Implementation Issue



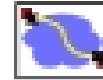
- Operating system timers are very coarse – how to pace packets out smoothly?
- Implemented using a congestion window that limits how much data can be in the network.
  - TCP also keeps track of how much data is in transit
- Data can only be sent when the amount of outstanding data is less than the congestion window.
  - The amount of outstanding data is increased on a "send" and decreased on "ack"
  - $(\text{last sent} - \text{last acked}) = \text{congestion window}$
- Window limited by both congestion and buffering
  - $\text{Sender's maximum window} = \text{Min}(\text{advertised window, send})$

11-0-07

Lecture 10: TCP Congestion Control

7

## ACK Clocking



- Congestion window helps to "pace" the transmission of data packets
- In steady state, a packet is sent when an ack is received
  - Data transmission remains smooth, once it is smooth
  - Self-clocking behavior



11-0-07

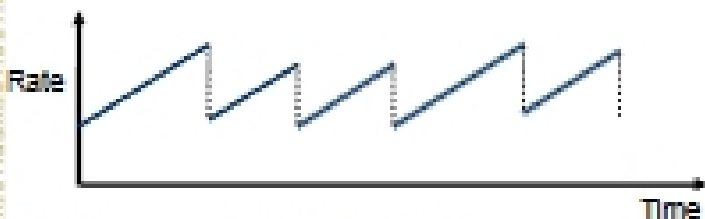
Lecture 10: TCP Congestion Control

8

## AIMD



- Distributed, fair and efficient
- Packet loss is seen as sign of congestion and results in a multiplicative rate decrease
  - Factor of 2
- TCP periodically probes for available bandwidth by increasing its rate



11-0-07

Lecture 10: TCP Congestion Control

9

## Congestion Avoidance



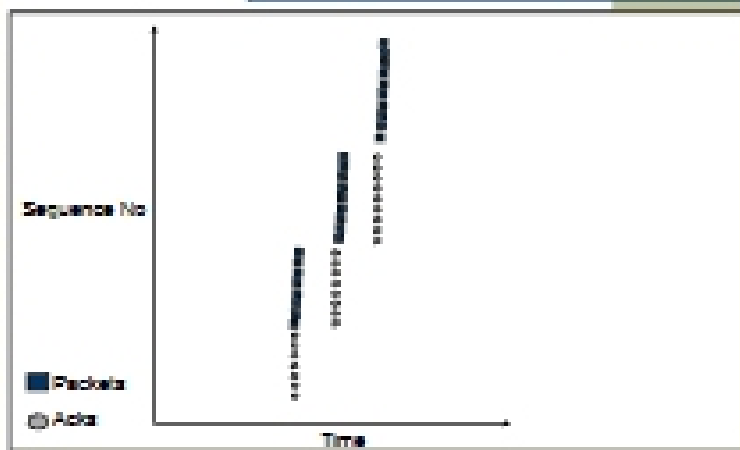
- If loss occurs when  $cwnd = W$ 
  - Network can handle  $0.5W \sim W$  segments
  - Set  $cwnd$  to  $0.5W$  (multiplicative decrease)
- Upon receiving ACK
  - Increase  $cwnd$  by  $(1 \text{ packet})/cwnd$ 
    - What is 1 packet?  $\rightarrow$  1 MSS worth of bytes
    - After  $cwnd$  packets have passed by  $\rightarrow$  approximately increase of 1 MSS
- Implements AIMD

11-0-07

Lecture 10: TCP Congestion Control

10

## Congestion Avoidance Sequence Plot

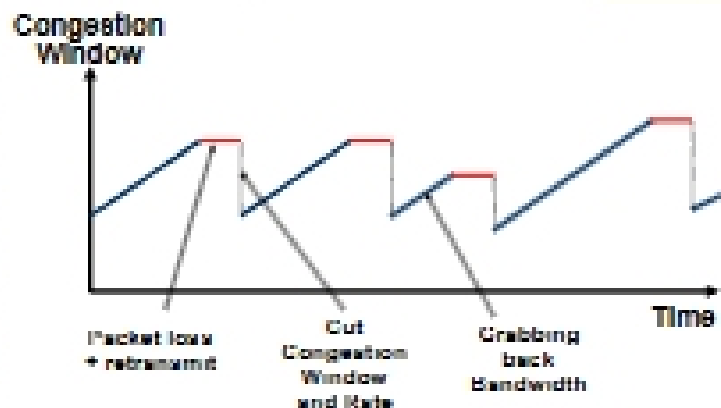


11-0-07

Lecture 10: TCP Congestion Control

11

## Congestion Avoidance Behavior



11-0-07

Lecture 10: TCP Congestion Control

12