

Reaching Definitions

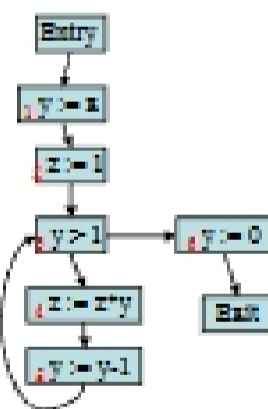
Done many ways

How many ways can we do Reaching Definitions?

- Data Flow analysis
- Constraint based approach
- Abstract Interpretation
- Type and Effect Systems

Running Example

```
y := x; // 1
z := 1; // 2
while y > 1 do { // 3
  z := z*y; // 4
  y := y-1; // 5
}
y := 0 // 6
```



Data Flow Analysis

- $RD_{\text{exit}}(1) = RD_{\text{entry}}(1) - \{ (y, l) \mid l \text{ in Lab } \} \cup \{ (y, 1) \}$
- $RD_{\text{exit}}(2) = RD_{\text{entry}}(2) - \{ (z, l) \mid l \text{ in Lab } \} \cup \{ (z, 2) \}$
- ...
- $RD_{\text{entry}}(1) = \{ (x, ?), (y, ?), (z, ?) \}$
- $RD_{\text{entry}}(2) = RD_{\text{exit}}(1)$
- $RD_{\text{entry}}(3) = RD_{\text{exit}}(2) \cup RD_{\text{exit}}(5)$
- ...

Finding a solution

- Let F be this set of equations
- Find RD s.t. $RD = F(RD)$
 - find least such solution
 - what is a non-minimal solution?

Constraint based approach

- $RD_{\text{exit}}(1) \supseteq RD_{\text{entry}}(1) - \{(y, l) \mid l \text{ in Lab}\}$
- $RD_{\text{exit}}(1) \supseteq \{(y, 1)\}$
- $RD_{\text{exit}}(2) \supseteq RD_{\text{entry}}(2) - \{(z, l) \mid l \text{ in Lab}\}$
- $RD_{\text{exit}}(2) \supseteq \{(z, 2)\}$
- ...
- $RD_{\text{entry}}(1) \supseteq \{(x, ?), (y, ?), (z, ?)\}$
- $RD_{\text{entry}}(2) \supseteq RD_{\text{exit}}(1)$
- $RD_{\text{entry}}(3) \supseteq RD_{\text{exit}}(2)$
- $RD_{\text{entry}}(3) \supseteq RD_{\text{exit}}(5)$
- ...

Something different

- $[[\text{fn } x \Rightarrow [x]^1]^2 [\text{fn } y \Rightarrow [y]^3]^4]^5$
- For each function application, which function may be applied?

- $C(l)$ = values that subexpression l can evaluate to
- $p(x)$ = values that x can be bound to
- $[[\text{fn } x \Rightarrow [x]^1]^2 [\text{fn } y \Rightarrow [y]^3]^4]^5$
- $\{ \text{fn } x \Rightarrow [x]^1 \} \subseteq C(2)$
- $\{ \text{fn } y \Rightarrow [y]^3 \} \subseteq C(4)$

Values taken on by a variable use

- $[[\text{fn } x \Rightarrow [x]^1]^2 [\text{fn } y \Rightarrow [y]^3]^4]^5$
- $p(x) \subseteq C(1)$
- $p(y) \subseteq C(3)$

Results of function application

- $[[\text{fn } x \Rightarrow [x]^1]^2 [\text{fn } y \Rightarrow [y]^3]^4]^5$
- Only one function application
 - but we don't know what function will be applied
- If function might be $\text{fn } x \Rightarrow [x]^1$
 - $\{ \text{fn } x \Rightarrow [x]^1 \} \subseteq C(2) \Rightarrow C(4) \subseteq p(x)$
 - $\{ \text{fn } x \Rightarrow [x]^1 \} \subseteq C(2) \Rightarrow C(1) \subseteq C(5)$
- If function might be $\text{fn } y \Rightarrow [y]^3$
 - $\{ \text{fn } y \Rightarrow [y]^3 \} \subseteq C(2) \Rightarrow C(4) \subseteq p(y)$
 - $\{ \text{fn } y \Rightarrow [y]^3 \} \subseteq C(2) \Rightarrow C(3) \subseteq C(5)$

Least solution

- $[[\text{fn } x \Rightarrow [x]^1]^2 [\text{fn } y \Rightarrow [y]^3]^4]^5$
- $C(1) = \{ \text{fn } y \Rightarrow [y]^3 \}$
- $C(2) = \{ \text{fn } x \Rightarrow [x]^1 \}$
- $C(3) = \{ \}$
- $C(4) = \{ \text{fn } y \Rightarrow [y]^3 \}$
- $C(5) = \{ \text{fn } y \Rightarrow [y]^3 \}$
- $p(x) = \{ \text{fn } y \Rightarrow [y]^3 \}$
- $p(y) = \{ \}$

Abstract Interpretation

- A trace is a series of (label, variable) pairs indicating the sequence of assignments
- A collecting semantics records the set of traces tr that can reach a given program point
- $RD(tr)(x) = l$ iff the rightmost pair (x, l') in tr has $l = l'$