

## Garbage Collection

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## What is garbage collection?

- Automatic recycling of allocated heap memory that can not be used again
  - i.e., is garbage
  - i.e., is not reachable (from roots)
- Convenient
- Avoids memory leaks (usually)
- Required for type safety

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## Java can leak memory

```
// Find the error
class Stack {
    Object a[] = new Object[10];
    int top = -1;
    Object pop() {
        if (top == -1)
            throw new NoSuchElementException();
        return a[top--];
    }
    void push(Object o) {
        if (top+1 == a.length)
            throw new StackOverflowException();
        a[++top] = o;
    }
}
```

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## Reachable != will be used again

```
// Error fixed in green
class Stack {
    Object a[] = new Object[10];
    int top = -1;
    Object pop() {
        if (top == -1)
            throw new NoSuchElementException();
        Object r = a[top];
        a[top--] = null;
        return r;
    }
    void push(Object o) {
        if (top+1 == a.length)
            throw new StackOverflowException();
        a[++top] = o;
    }
}
```

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## Finding pointers can be hard

- What are root pointers?
  - Pointers that can be accessed by program, even though no pointers point to them
  - For Java:
    - all local and stack variables of all threads
    - all loaded classes (class unloading?)
    - static variables of loaded classes
    - What is the type of a Java stack variable?
      - Need stack map, from PC to local/stack variable types
      - Can also take into account dead references

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## Finding pointers in objects can be hard

- Type unsafe languages are a nightmare
  - variant records (unions)
  - pointers into middle of objects
  - various bit-twiddling tricks
  - Even in typesafe languages
    - need run-time access to type information for objects

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## Reference counting

- Different approach to garbage collection
- For each object, keep track of number of references to it
- Cost to adjust counts (some optimization possible)
  - {Object tmp = a; a = b; b = tmp}
  - generates:  
a.refCount++; tmp = a;  
b.refCount++; a.refCount--; a = b;  
tmp.refCount++; b.refCount--; b = tmp;  
tmp.refCount--;
- Can't reclaim circular structures
- Use only in special circumstances
  - not generally recommended

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## Mark and Sweep [McCarthy 1960]

- From roots
  - Roots are pointers known to be accessible
    - e.g., registers, stack
- perform a depth-first search to mark live nodes
- Sweep through all memory
  - if node is unmarked, it is garbage, put on free list
  - if marked, is in use, unmark
    - to prepare it for next garbage collection

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## Problems with Mark and Sweep

- Many of these are problems with other garbage collection schemes
  - Stack required for DFS
    - could be big
  - Finding roots
  - Finding pointers
    - Easy for CONS cells
  - Stops the world

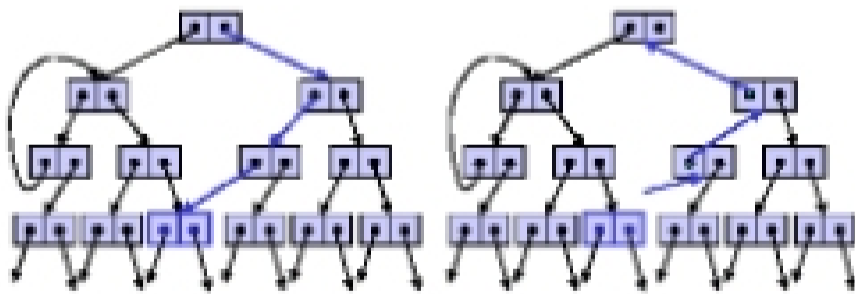
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## Pointer Reversal

- In depth first search, need stack as big as depth of data structure
  - for a long list, depth == length
- Often, garbage is being collected when we are short on space
- Can avoid storing DFS stack separately using pointer reversal
  - Efficient way to implement mark and sweep

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## Doing pointer reversal



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## Is pointer reversal worthwhile?

- Requires additional bits
  - $\log f$  bits, where  $f$  is number of pointer fields
- Requires additional visits and modifications of each node
- Generally not worth while
  - Instead, can do tail-call optimization for last field from object
  - when following last field out of an object, don't put object on stack for a return visit
  - helps a little

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