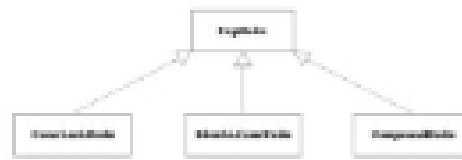


Inheritance Hierarchies

- Because expressions have more than one form, a C++ class that represents expressions can be represented most easily by a class hierarchy in which each of the expression types is a separate subclass, as shown in the following diagram:



- Even though the class hierarchy is organized in terms of the different types of nodes, clients of the expression package will almost always work with pointers to nodes instead. As I did last time, I will therefore give the pointer type the name `expressionT`.

Representing Inheritance in C++

- The first step in creating a C++ subclass is to indicate the superclass on the header line, using the following syntax:

```
class subclass: public superclass {
    body of class definition
}
```

- In contrast to Java, a subclass cannot automatically override the definition of a method in its superclass. To permit such overriding, both classes must mark the prototype for that method with the keyword `virtual`.
- An **abstract class** is a class that doesn't actually represent any objects but instead serves only as a common superclass for **concrete classes** that do generate objects. In C++, methods for an abstract class that are always implemented by the concrete subclasses are indicated by including `= 0` before the semicolon on the prototype line.

The `exp.h` Interface

```

/*
 * file: exp.h
 * -----
 * This interface defines a class hierarchy for expressions,
 * which allows the client to represent and manipulate algebraic
 * binary expression trees.
 */

#ifndef _exp_h
#define _exp_h

#include "symbol.h"
#include "exp.h"

/*
 * type: expression
 * -----
 * For clients, the main interface type supported by this interface
 * is the expressionT type, which is defined as a pointer to an
 * expression object. This is the type used by all other functions
 * and methods in the expression package.
 */

class expression;
typedef expression *expressionT;
```

The `exp.h` Interface

```

/*
 * class: evaluation
 * -----
 * This class is passed by reference through the recursive levels
 * of the evaluation and contains information from the evaluation
 * environment when the evaluator may need to look. The only
 * such information implemented here is a symbol table when the
 * variable names have their values.
 */

class evaluation {
public:
    evaluation();
    evaluation(int);
    void evaluate(symbol var, int value);
    int evaluate(symbol var);
    bool isdefined(symbol var);
private:
    symbolTable symbolTable;
};
```

The `exp.h` Interface

```

/*
 * class: symbol
 * -----
 * This class is used to represent a node in an expression tree.
 * symbol is an example of an abstract class, which defines the
 * structure and behavior of a set of classes but has no objects
 * of its own. Any objects must be one of the three concrete
 * subclasses of symbol:
 *
 * 1. constant -- an integer constant
 * 2. identifier -- a string representing an identifier
 * 3. unary -- one expression combined by an operator
 *
 * The symbol class defines the interface common to all symbol
 * objects: each subclass provides its own specific implementation
 * of the common interface.
 *
 * Note on syntax: each of the virtual methods in the symbol class
 * is marked with the designation = 0 on the prototype line. This
 * notation is used in all to indicate that this method is purely
 * virtual and will always be supplied by the subclass.
 */
```

The `exp.h` Interface

```

/* expression type for the three expression types */
enum expType { constantType, identifierType, unaryType };

/* constant array for symbol class */

class symbol {
public:
    symbol();
    virtual ~symbol();
    virtual expType type() const = 0;
    virtual int val(evaluation & env) const = 0;
    virtual string name() const = 0;
};
```

The exp.h Interface

```

/*
 * class: constantnode
 * -----
 * this subclass represents a constant integer expression.
 */
class constantnode: public expnode {
public:
    constantnode(int val):
        virtual expnode::expnode() {
        virtual int eval(constenv & env):
        virtual string toString();
};
/*
 * method: generate
 * usage: value = ((constantnode *) exp)-generate();
 * -----
 * this method returns the value field without calling eval.
 */
int generate();
private:
    int value;
};

```

The exp.h Interface

```

/*
 * class: identifiernode
 * -----
 * this subclass represents an expression corresponding to a variable.
 */
class identifiernode: public expnode {
public:
    identifiernode(string name):
        virtual expnode::expnode() {
        virtual int eval(constenv & env):
        virtual string toString();
};
/*
 * method: generate
 * usage: name = ((identifiernode *) exp)-generate();
 * -----
 * this method returns the name field of the identifier node.
 */
string generate();
private:
    string name;
};

```

The exp.h Interface

```

/*
 * class: compoundnode
 * -----
 * this subclass represents a compound expression consisting of
 * two subexpressions joined by an operator.
 */
class compoundnode: public expnode {
public:
    compoundnode(char op, expnode *l, expnode *r):
        virtual compoundnode() {
        virtual expnode::expnode() {
        virtual int eval(constenv & env):
        virtual string toString();

    char op;
    expnode *l;
    expnode *r;
};
private:
    char op;
    expnode *l;
    expnode *r;
};

```

Code for the eval Method

```

int constantnode::eval(constenv & env) {
    return value;
}

int identifiernode::eval(constenv & env) {
    if (!env.lookup(name)) error("illegal identifier");
    return env.generate(name);
}

int compoundnode::eval(constenv & env) {
    if (op == '+') {
        int lval = ((identifiernode *) l)-eval(env);
        int rval = ((identifiernode *) r)-eval(env);
        return lval + rval;
    }
    if (op == '-') {
        int lval = ((identifiernode *) l)-eval(env);
        int rval = ((identifiernode *) r)-eval(env);
        return lval - rval;
    }
    if (op == '*') {
        int lval = ((identifiernode *) l)-eval(env);
        int rval = ((identifiernode *) r)-eval(env);
        return lval * rval;
    }
    if (op == '/') {
        int lval = ((identifiernode *) l)-eval(env);
        int rval = ((identifiernode *) r)-eval(env);
        return lval / rval;
    }
    error("illegal operator in expression");
}

```

A Two-Level Grammar

- The problem of parsing an expression can be simplified by changing the grammar to one that has two levels:
 - An **expression** is either a **term** or two expressions joined by an operator.
 - A **term** is either a constant, an identifier, or an expression enclosed in parentheses.
- This design is reflected in the following revised grammar.

```

E → T
E → E op E

T → constant
T → identifier
T → ( E )

```

The parser.cpp Implementation

```

/*
 * implementation name: make
 * usage: exp = make(scanner, gram);
 * -----
 * this function reads the raw expression from the scanner by
 * making the input to the following ambiguous grammar:
 *
 *   E → T
 *   E → E op E
 *
 * this version of the method uses precedence to resolve ambiguity.
 */
expnode* make(scanner & scanner, int gram) {
    expnode* exp = make(scanner);
    string value;
    while (true) {
        value = scanner.nexttoken();
        int success = evaluate(value);
        if (success == gram) break;
        expnode* rhs = make(scanner, success);
        exp = new compoundnode(value[0], exp, rhs);
    }
    scanner.advance(value);
    return exp;
}

```