


CSE 3302  
Programming Languages




## Syntax (cont.)

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Spring 2008

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## What is Parsing?



- Given a grammar and a token string:
  - determine if the grammar can generate the token string?
  - i.e., is the string a legal program in the language?
- In other words, to construct a parse tree for the token string.

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## What's significant about parse tree?




A parse tree gives a unique syntactic structure

- Leftmost, rightmost derivation
- There is only one leftmost derivation for a parse tree, and symmetrically only one rightmost derivation for a parse tree.

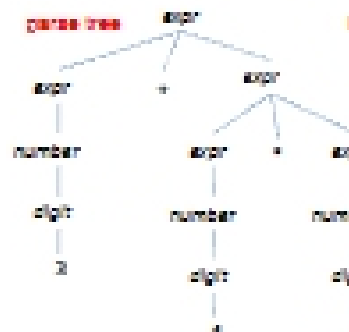
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## Example




$expr \rightarrow expr + expr \mid expr * expr \mid ( expr ) \mid number$   
 $number \rightarrow number digit \mid digit$   
 $digit \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$

parse tree




leftmost derivation




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## What's significant about parse tree?




A parse tree has a unique meaning, thus provides basis for semantic analysis.  
(Syntax-directed semantics: semantics are attached to syntactic structure.)



$expr.val = expr1.val + expr2.val$

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## Relationship among language, grammar, parser



- Chomsky Hierarchy  
[https://en.wikipedia.org/wiki/Chomsky\\_hierarchy](https://en.wikipedia.org/wiki/Chomsky_hierarchy)
- A language can be described by multiple grammars, while a grammar defines one language.
- A grammar can be parsed by multiple parsers, while a parser accepts one grammar, thus one language.
- Should design a language that allows simple grammar and efficient parser
- For a language, we should construct a grammar that allows fast parsing
- Given a grammar, we should build an efficient parser

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## Ambiguity

- **Ambiguous grammar:** There can be multiple parse trees for the same sentence (program)
  - In other words, multiple leftmost derivations.
- **Why is it bad?**
  - Multiple meanings

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## Ambiguity

- **Was this ambiguous?**

```
number → number digit | digit
digit → 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9
```
- **How about this?**

```
expr → expr - expr | number
```

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## Deal with Ambiguity

- **Unambiguous Grammar**
  - Rewrite the grammar to avoid ambiguity.

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## Example of Ambiguity: Precedence

```
expr → expr + expr | expr * expr | ( expr ) | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9
```

Two different parse trees for expression 3+4\*5

parse tree 1

parse tree 2

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## Eliminating Ambiguity for Precedence

- Establish “precedence cascade”: using different structured names for different constructs, adding grammar rules.
  - Higher precedence : lower in cascade

```
expr → expr + expr | expr * expr | ( expr ) | number
```

➔

```
expr → expr + expr | term
term → term * term | ( expr ) | number
```

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## Example of Ambiguity: Associativity

```
expr → expr - expr | ( expr ) | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9
```

Two different parse trees for expression 5-2-1

parse tree 1

expr.val = 4

parse tree 2


expr.val = 2

– is right-associative, which is against common practice in integer arithmetic

– is left-associative, which is what we usually assume

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## Associativity




- Left-Associative: + - \* /
- Right-Associative: =

What is meant by a=b=c=1?

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## Eliminating Ambiguity for Associativity



- **left-associativity: left-recursion**

```

expr → expr - expr | ( expr ) | number
      ↓
expr → expr - term | term
term → (expr) | number
    
```


- **right-associativity: right-recursion**

```

expr → expr = expr | a | b | c
      ↓
expr → term = expr | term
term → a | b | c
    
```

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## Putting Together



```

expr → expr - expr | expr / expr | ( expr ) | number
number → number digit | digit
digit → 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9
    
```

We want to make: - left-associative and / has precedence over -


→

```

expr → expr - term | term
term → term / factor | factor
factor → ( expr ) | number
number → number digit | digit
digit → 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9
    
```

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## Example of Ambiguity: Dangling-Else




```


stmt → !E ( expr ) stmt
      | !E ( expr ) stmt else stmt
      | other-stmt
    
```

Two different parse trees for "!E ( expr ) !E ( expr ) stmt else stmt"

parse tree 1




parse tree 2



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## Eliminating Dangling-Else




```

stmt → matched_stmt
      | unmatched_stmt
matched_stmt → !E ( expr ) matched_stmt else matched_stmt
              | other_stmt
unmatched_stmt → !E ( expr ) stmt
                | !E ( expr ) matched_stmt else unmatched_stmt
    
```

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## EBNF



- Repetition { }

```

number → digit | number digit   → number → { digit }
expr → expr - term | term       → expr → term [ - term ]
    
```

- Option [ ]

```

signed-number → sign number / number   → signed-number → [ sign ] number
!-stmt → !E ( expr ) stmt              → !-stmt → !E ( expr ) stmt [ else stmt ]
    
```

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