

## CSE 321 Discrete Structures

Winter 2008  
Lecture 1  
Propositional Logic

## About the course

- From the CSE catalog:
  - CSE 321 Discrete Structures (4)  
Fundamentals of set theory, graph theory, enumeration, and algebraic structures, with applications in computing. Prerequisite: CSE 143; either MATH 126, MATH 129, or MATH 136.
- What I think the course is about:
  - Foundational structures for the practice of computer science and engineering

## Why this material is important

- Language and formalism for expressing ideas in computing
- Fundamental tasks in computing
  - Translating imprecise specification into a working system
  - Getting the details right

## Topic List

- Logic/boolean algebra: hardware design, testing, artificial intelligence, software engineering
- Mathematical reasoning/induction: algorithm design, programming languages
- Number theory/probability: cryptography, security, algorithm design, machine learning
- Relations/relational algebra: databases
- Graph theory: networking, social networks, optimization

## Administration

- Instructor
  - Richard Anderson
- Teaching Assistant
  - Natalie Linnell
- Quiz section
  - Thursday, 12:30 – 1:20, or 1:30 – 2:20
  - CSE 305
- Recorded Lectures
  - Available on line
- Text: Rosen, Discrete Mathematics
  - 6<sup>th</sup> Edition preferred
  - 5<sup>th</sup> Edition okay
- Homework
  - Due Wednesdays (starting Jan 16)
- Exams
  - Midterms, Feb 8
  - Final, March 17, 2:30-4:20 pm
- All course information posted on the web
- Sign up for the course mailing list

## Propositional Logic

## Propositions

- A statement that has a truth value
- Which of the following are propositions?
  - The Washington State flag is red
  - It snowed in Whistler, BC on January 4, 2008.
  - Hillary Clinton won the democratic caucus in Iowa
  - Space aliens landed in Roswell, New Mexico
  - Ron Paul would be a great president
  - Turn your homework in on Wednesday
  - Why are we taking this class?
  - If  $n$  is an integer greater than two, then the equation  $a^n + b^n = c^n$  has no solutions in non-zero integers  $a$ ,  $b$ , and  $c$ .
  - Every even integer greater than two can be written as the sum of two primes
  - This statement is false
- Propositional variables:  $p, q, r, s, \dots$
- Truth values: **T** for true, **F** for false

## Compound Propositions

- Negation (not)  $\neg p$
- Conjunction (and)  $p \wedge q$
- Disjunction (or)  $p \vee q$
- Exclusive or  $p \oplus q$
- Implication  $p \rightarrow q$
- Biconditional  $p \leftrightarrow q$

## Truth Tables

$p$	$\neg p$

$p$	$q$	$p \wedge q$

$p$	$q$	$p \vee q$

$p$	$q$	$p \oplus q$

## Understanding complex propositions

- Either Harry finds the locket and Ron breaks his wand or Fred will not open a joke shop

## Understanding complex propositions with a truth table

$p$	$r$	$f$	$p \wedge r$	$\neg f$	$(p \wedge r) \oplus \neg f$

## Aside: Number of binary operators

- How many different binary operators are there on atomic propositions?

$$p \rightarrow q$$

$p$	$q$	$p \rightarrow q$

- Implication
  - $p$  implies  $q$
  - whenever  $p$  is true  $q$  must be true
  - if  $p$  then  $q$
  - $q$  if  $p$
  - $p$  is sufficient for  $q$
  - $p$  only if  $q$

If pigs can whistle then horses can fly

### Converse, Contrapositive, Inverse

- Implication:  $p \rightarrow q$
- Converse:  $q \rightarrow p$
- Contrapositive:  $\neg q \rightarrow \neg p$
- Inverse:  $\neg p \rightarrow \neg q$
- Are these the same?

### Biconditional $p \leftrightarrow q$

- $p$  iff  $q$
- $p$  is equivalent to  $q$
- $p$  implies  $q$  and  $q$  implies  $p$

$p$	$q$	$p \leftrightarrow q$

### English and Logic

- You cannot ride the roller coaster if you are under 4 feet tall unless you are older than 16 years old
  - $q$ : you can ride the roller coaster
  - $r$ : you are under 4 feet tall
  - $s$ : you are older than 16

### Logical equivalence

- Terminology: A compound proposition is a
  - Tautology if it is always true
  - Contradiction if it is always false
  - Contingency if it can be either true or false

$$p \vee \neg p$$

$$(p \oplus p) \vee p$$

$$p \oplus \neg p \oplus q \oplus \neg q$$

$$(p \rightarrow q) \wedge p$$

$$(p \wedge q) \vee (p \wedge \neg q) \vee (\neg p \wedge q) \vee (\neg p \wedge \neg q)$$