

Dynamic Memory Allocation II

April 1, 2004

Topics

- Explicit doubly-linked free lists
- Segregated free lists
- Garbage collection
- Memory-related perils and pitfalls

Freeing With Explicit Free Lists

Insertion policy: Where in the free list do you put a newly freed block?

- LIFO (last-in-first-out) policy
 - Insert freed block at the beginning of the free list
 - Proc: simple and constant time
 - Conc: studies suggest fragmentation is worse than address ordered.
- Address-ordered policy
 - Insert freed blocks so that free list blocks are always in address order
 - I.e. $\text{addr}(\text{prev}) < \text{addr}(\text{curr}) < \text{addr}(\text{succ})$
 - Conc: requires search
 - Proc: studies suggest fragmentation is better than LIFO

Freeing With a LIFO Policy

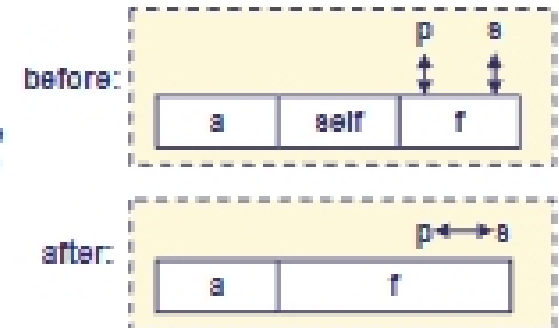
Case 1: a-a-a

- Insert self at beginning of free list



Case 2: a-a-f

- Splice out next, coalesce self and next, and add to beginning of free list



Freeing With a LIFO Policy

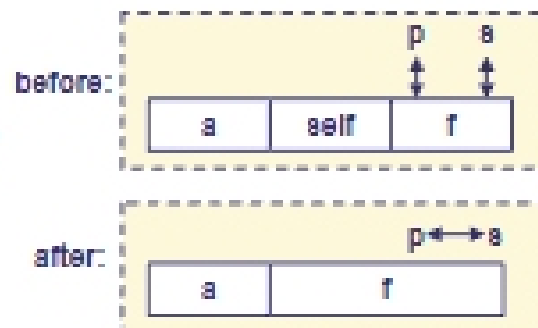
Case 1: a-a-a

- Insert self at beginning of free list



Case 2: a-a-f

- Splice out next, coalesce self and next, and add to beginning of free list



Freeing With a LIFO Policy (cont)

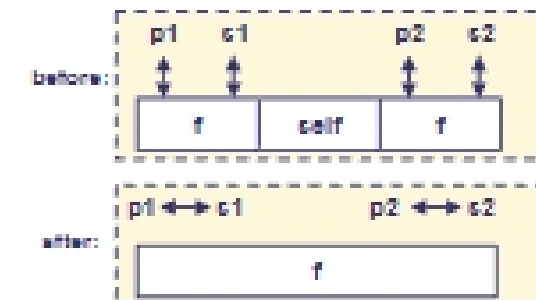
Case 3: f-a-a

- Splice out prev, coalesce with self, and add to beginning of free list



Case 4: f-a-f

- Splice out prev and next, coalesce with self, and add to beginning of list



Explicit List Summary

Comparison to implicit list:

- Allocate is linear time in number of free blocks instead of total blocks -- much faster allocates when most of the memory is full
- Slightly more complicated allocate and free since needs to splice blocks in and out of the list
- Some extra space for the links (2 extra words needed for each block) Does this increase internal frag?

Main use of linked lists is in conjunction with segregated free lists

- Keep multiple linked lists of different size classes, or possibly for different types of objects

Keeping Track of Free Blocks

Method 1: *Implicit list* using lengths -- links all blocks



Method 2: *Explicit list* among the free blocks using pointers within the free blocks

Method 3: *Segregated free list*

- Different free lists for different size classes

Method 4: Blocks sorted by size

- Can use a balanced tree (e.g. Red-Black tree) with pointers within each free block, and the length used as a key