

Summary From the Last Lecture

- ◆ Key exchange
 - Kerberos
 - Digital certificates
- ◆ Certificate authority structure
 - PGP, hierarchical model
- ◆ Recovery from exposed keys
 - Revocation lists, time-limited keys, real time validation
- ◆ Group key management
 - Robustness, forward/backward secrecy

Authentication

Basis for Authentication

- ◆ Ideally
 - Who you are
- ◆ Practically
 - Something you know
 - Something you have
 - Something about you

Something You Know

- ◆ Password or Algorithm
 - e.g. encryption key derived from password
- ◆ Issues
 - Someone else may learn it
 - Find it, sniff it, trick you into providing it
 - Other party must know how to check
 - You must remember it

Password Authentication

- ◆ Alice inputs her password, computer verifies this against list of passwords
- ◆ If computer is broken into, hackers can learn everybody's passwords
 - Use one-way functions, store the result for every valid password
 - Perform one-way function on input, compare result against the list

Password Authentication

- ◆ Hackers can compile a list of frequently used passwords, apply one-way function to each and store them in a table - **dictionary attack**
- ◆ Host adds random salt to password, applies one-way function to that and stores result and salt value
 - Randomly generated, unique and long enough

Password Authentication

- ◆ Someone sniffing on the network can learn the password
- ◆ Lamport hash or S-KEY – time-varying pass
 - To set-up the system, Alice enters random number R
 - Host calculates $x_0=h(R), x_1=h(h(R)), x_2=h(h(h(R))), \dots, x_{100}$
 - Alice keeps this list, host sets her password to x_{100}
 - Alice logs on with x_{100} , host verifies $h(x_{100})=x_{101}$, resets password to x_{100}
 - Next time Alice logs on with x_{99}

Password Authentication

- ◆ Someone sniffing on the network can learn the password
 - Host keeps a file of every user's public key
 - Users keep their private keys
 - When Alice attempts to log on, host sends her a random number R
 - Alice encrypts R with her private key and sends to host
 - Host can now verify her identity by decrypting the message and retrieving R

Public Key Authentication

- ◆ Key Distribution
 - Confidentiality not needed for public key
 - Solves n^2 problem
- ◆ Performance
 - Slower than conventional cryptography
 - Implementations used for key distribution, then use conventional crypto for data encryption
- ◆ Trusted third party still needed
 - To certify public key
 - To manage revocation
 - In some cases, third party may be off-line

Single Sign-On

- ◆ Passport
- ◆ Liberty Alliance
- ◆ Shibboleth

Federated Identity Passport vs Liberty Alliance

- ◆ Two versions of Passport
 - Centralized and federated
- ◆ Liberty Alliance
 - Loosely federated with framework to describe authentication provided by others

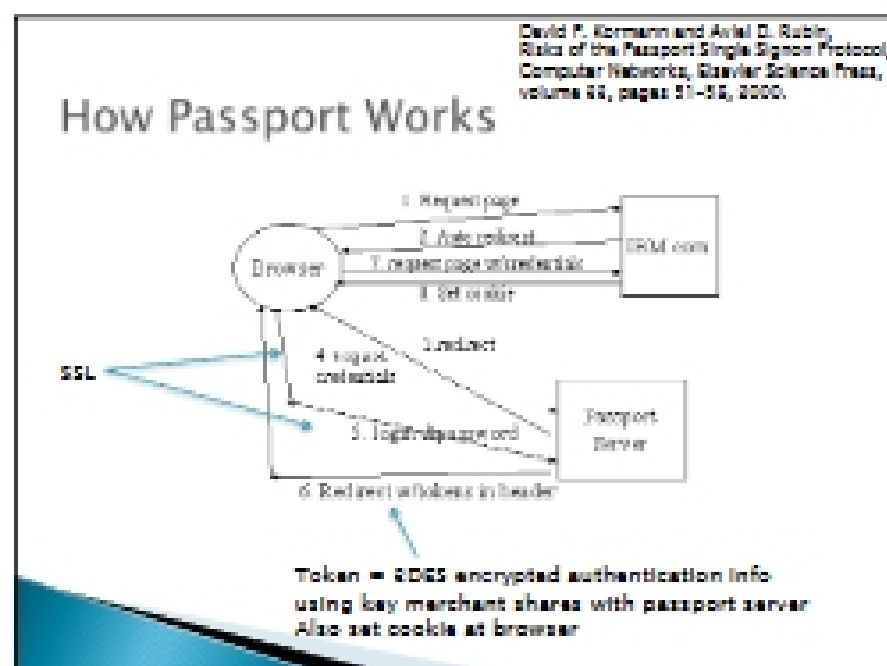
Passport v1

- ◆ Goal is single sign-on
 - Solves problem of weak or repeated user/pass combinations
- ◆ Implemented via redirections
 - Users authenticate themselves to a common server, which gives them tickets
 - Similar flavor to Kerberos but different environment – many organizations
- ◆ Widely deployed by Microsoft
 - Designed to use existing technologies in servers/browsers (HTTP redirect, SSL, cookies, Javascript)

David P. Kormann and Aviel D. Rubin,
Risk of the Passport Single Signon Protocol,
Computer Networks, Elsevier Science Press,
volume 34, pages 21-36, 2000.

How Passport Works

- ◆ Client (browser), merchant (Web server), Passport login server
- ◆ Passport server maintains authentication info for client
 - Gives merchant access when permitted by client
- ◆ Divides client data into profile (address) and wallet (credit card)



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Some Problems with Passport

- ◆ User interface is confusing and may misrepresent the reality
- ◆ Weak keys may be used for 3DES
- ◆ Single key is used to encrypt cookies for all clients
- ◆ Cookies stay on machine, can be stolen
 - No authenticator (timestamp), like in Kerberos, enables reuse by others
- ◆ Coupling of Hotmail with Passport

Read more at <http://avtrubin.com/passport.html>

Federated Passport

- ◆ Multiple federated identity providers
 - E.g. ISPs register own users
 - One can rely on claims made by other ID providers
- ◆ Claims
 - Emails, relationships, authorization for scenarios, ownership of private/public key pair
- ◆ Need "translators" for different claim languages

Liberty Alliance

- ◆ Design criteria was most of the issues addressed by Federated Passport, i.e. no central authority
- ◆ Use SAML (Security Association Markup Language) to describe trust across authorities, and what assertions mean from particular authorities
- ◆ Four assurance levels
 - How much we trust a given identity assertion
 - Little, some, high and very high confidence