

Plan for Today

Recursive descent or predictive parsing

- example predictive parser
- FIRST and FOLLOW sets revisited
- constructing a predictive parser table

Syntax-directed translation

Example Predictive Parser

```
S -> Mesh EOF
Mesh -> num Modelist num ElemList
Modelist -> ε | Node Modelist
Node -> node num real real // node_id, x, y
ElemList -> ε | Elem ElemList
Elem -> tri num num num num // elem_id, 3 node ids
Elem -> sqr num num num num num // elem_id, 4 node ids
```

```
void S() { switch(tok) {
  case NUM: Mesh(); set(EOF); break;
  default: error();
}}

void Mesh() { switch(tok) {
  case NUM: num_nodes = NUM.val; set(NUM);
             Modelist();
             num_elem = NUM.val; set(EOF); break;
  default: error();
}}

void Modelist() { switch(tok) {
  case NUM: break;
  case NODE: Node(); Modelist(); break;
  default: error();
}}

void Node() { switch(tok) {
  case NUM: num = NUM.val; set(NUM);
             real1 = REAL.val; set(REAL);
             real2 = REAL.val; set(REAL); break;
  default: error();
}}

void ElemList() { switch(tok) {
  case TRIP: Elem(); set(EOF); break;
  case SQR: Elem(); set(EOF); break;
  default: error();
}}

void Elem() { switch(tok) {
  case TRIP: elem_id = NUM.val; set(NUM);
             num1 = NUM.val; set(NUM);
             num2 = NUM.val; set(NUM);
             num3 = NUM.val; set(NUM); break;
  case SQR: elem_id = NUM.val; set(NUM);
             num1 = NUM.val; set(NUM);
             num2 = NUM.val; set(NUM);
             num3 = NUM.val; set(NUM);
             num4 = NUM.val; set(NUM); break;
  default: error();
}}

void error() { printf("Error: %s\n", tok); }
```

FIRST and FOLLOW sets

nullable(X)

- X is a nonterminal
- nullable(X) is true if X can derive the empty string

FIRST

- FIRST(z) = {z}, where z is a terminal
- FIRST(X) = \cup FIRST(rhs_i), where X is a nonterminal and $X \rightarrow$ rhs_i
 - union all of FIRST(sym) on rhs up to and including first nonnullable

FOLLOW(Y), only relevant when Y is a nonterminal

- look for Y in rhs of rules (lhs \rightarrow rhs) and union all FIRST sets for symbols after Y up to and including first nonnullable
- if all symbols after Y are nullable then also union in FOLLOW(lhs)

Constructing the Predictive Parser Table

Algorithm

```
for each X -> gamma
  for each T in FIRST(gamma)
    table[X,T] = X->gamma
  if gamma is nullable
    for each T in FOLLOW(X)
      table[X,T] = X->gamma
```

```
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```

Syntax-directed translation

One-pass versus multi-pass compiler

- What do we need for MiniJava?

In a predictive, or top-down, parser

- do actions in functions for nonterminals
- example: storing off the number of nodes and elements in the mesh grammar parser

In a shift-reduce, or bottom-up, parser

- each reduction leads to an action being performed
- example: SableCC builds a parse tree using actions associated with each grammar rule
- syntax-directed translation is equivalent to performing an action on internal nodes in the parse tree during an in post-order traversal
- example: interpreter you wrote for PA1